



## **Residue Number Systems**



# A Puzzle

What number has the reminders 2, 3 and 2  
when divided by the numbers 7, 5 and 3 ?

$$x \bmod 7 = 2$$

$$x \bmod 5 = 3$$

$$x \bmod 3 = 2$$

$$x = ???$$

Chinese mathematician Sun Tzu - third-century AD  
(not the famous military strategist Sun Tzu)



# Chinese Remainder Theorem

There exists an integer  $x$  solving the system of equations:

$$x \bmod m_1 = x_1$$

...

$$x \bmod m_n = x_n$$

where numbers  $m_1 \dots m_n$  are relatively prime.

# Residue Representation

## ► Convert X-value into RNS representation:

- find a set of numbers (i.e. moduli), which are all **relatively prime** to each other (greatest common divisor is 1):

$$m_{k-1}, m_{k-2}, \dots, m_1, m_0$$

$$m_{k-1} > m_{k-2} > \dots > m_1 > m_0$$

e.g. 7, 5, 3

- compute a set of residues with respect to moduli:

$$x_{k-1}, x_{k-2}, \dots, x_1, x_0 \quad \text{where } x_i = X \bmod m_i = \langle x \rangle_{m_i}$$

e.g. (X=12)  $\rightarrow$  5, 2, 0

- residue list is treated as a k-digit RNS number:

$$X = (x_{k-1} | x_{k-2} | \dots | x_1 | x_0)_{\text{RNS}(m_{k-1} | m_{k-2} | \dots | m_1 | m_0)} \quad \text{e.g. } (5|2|0)_{\text{RNS}(7|5|3)}$$

# RNS Examples

► Assume RNS(8|7|5|3)

$$123 = (3|4|3|0)_{\text{RNS}}$$

123				
moduli	8	7	5	3
residues	3	4	3	0

$$5 = (5|5|0|2)_{\text{RNS}}$$

5				
moduli	8	7	5	3
residues	5	5	0	2

$$0 = (0|0|0|0)_{\text{RNS}}$$

0				
moduli	8	7	5	3
residues	0	0	0	0



# Congruency Relation

- ▶ Two integers  $a$  and  $b$  are said to be congruent modulo  $n$ , if their difference  $(a-b)$  is an integer multiple of  $n$ .

$$a \equiv b \pmod{n}$$

- ▶  $X \equiv x \pmod{m} \rightarrow x = X - km$

e.g.

$$13 \equiv 3 \pmod{5} \rightarrow 3 = 13 - 2 \cdot 5$$

$$13 \equiv 8 \pmod{5} \rightarrow 8 = 13 - 1 \cdot 5$$

$$13 \equiv 13 \pmod{5} \rightarrow 13 = 13 - 0 \cdot 5$$

# Operations on Congruent Numbers

For congruencies  $A_i \equiv a_i$  (with modulus  $n$ ):

▶  $\sum A_i \equiv \sum a_i$

▶  $A_i - A_j \equiv a_i - a_j$

▶  $\prod A_i \equiv \prod a_i$

▶  $A_i^s \equiv a_i^s$

▶  $v \cdot A_i \equiv v \cdot a_i$

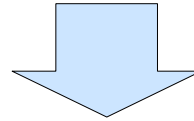
(except division)

	modulus = 5	
A1 ≡ a1	13	3
A2 ≡ a2	9	4
sum	22	7 (2)
diff.	4	-1 (4)
mult.	117	12 (2)

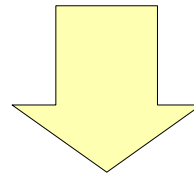


# Operations on RNS Numbers

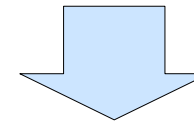
$$X = (x_{k-1} | x_{k-2} | \dots | x_1 | x_0)_{\text{RNS}(m_{k-1} | m_{k-2} | \dots | m_1 | m_0)}$$



$$X \equiv (x_{k-1} | x_{k-2} | \dots | x_1 | x_0)$$



$X \text{ oper } Y$



$$(x_{k-1} \text{ oper } y_{k-1} | \dots | x_1 \text{ oper } y_1 | x_0 \text{ oper } y_0)$$





# ***Operations on RNS Numbers***

- ▶ Operation (+, -, \*, ^) on RNS numbers is performed on all corresponding residues, totally in parallel
- ▶ Operation on residues at  $i$ -positions is performed modulo  $m_i$
- ▶ The resulting RNS number uniquely identifies the result of operation (within the dynamic range)



# Example of RNS Operations

- ▶ All operations are performed in parallel on corresponding residues (modulo  $m_i$ )

RNS (8|7|5|3)

A 21 (5|0|1|0)  
B 8 (0|1|3|2)

A+B 29 (5|1|4|2)  
A-B 13 (5|6|3|1)  
A\*B 168 (0|0|3|0)  
A<sup>2</sup> 441 (1|0|1|0)  
B<sup>3</sup> 512 (0|1|2|2)  
3\*A 63 (7|0|3|0)



# Representation Range

- ▶ RNS numbers may uniquely identify  $M$  numbers, where  $M = (m_{k-1} * \dots * m_1 * m_0) = \prod m_i$

e.g. RNS(8|7|5|3)  $\rightarrow 8*7*5*3 = 840$  unique combinations

- ▶  $M$  is called a dynamic range for a given RNS

- ▶ The range can cover any interval of  $M$ -consecutive numbers

e.g. for  $M=840$

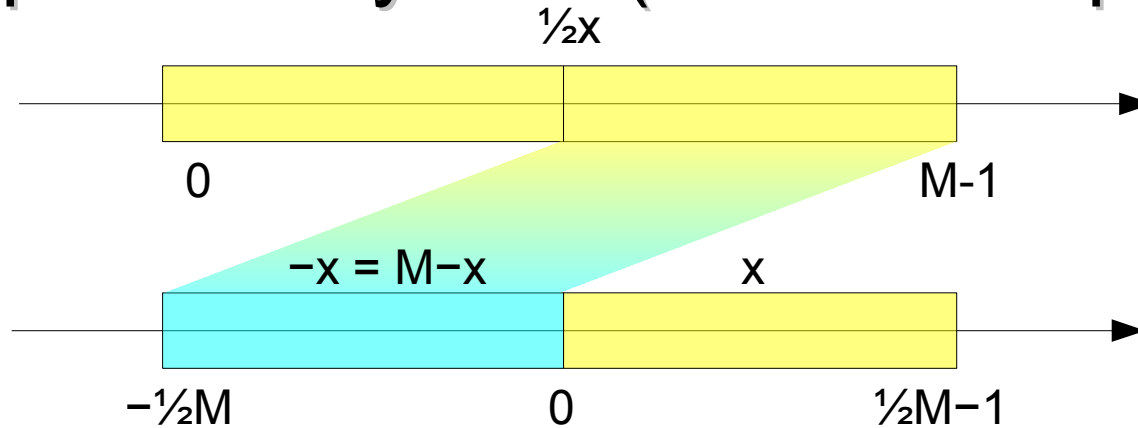
0..839, -420..419, etc.

$(0|0|0|0)_{\text{RNS}} = 0$  or 840 or 1680 ...

$(7|6|4|2)_{\text{RNS}} = 839$  or -1

# Negative RNS Numbers

- ▶ Negative numbers can be represented using a complement system (with  $M$ -complement)



- ▶ Residues of  $-x$  are equal to residues of  $M-x$

$$(-x) \bmod m_i = (M-x) \bmod m_i$$

- ▶ Residues of  $-x$  are  $m_i$ -complements of  $x$

$$x = (x_{k-1} \mid \dots \mid x_0) \rightarrow -x = (m_{k-1} - x_{k-1} \mid \dots \mid m_0 - x_0)$$

# Negative RNS - Examples

- ▶ Take RNS representation of positive number

$$x = 21 = (5|0|1|0)_{\text{RNS}} \quad (\text{RNS}(8|7|5|3))$$

- ▶ Calculate  $m_i$ -complements of all residues

$$(8-5 | 7-0 | 5-1 | 3-0)_{\text{RNS}} = (3 | 0 | 4 | 0)_{\text{RNS}} = -21$$

$$0 = (0|0|0|0)_{\text{RNS}} \rightarrow -0 = (0|0|0|0)_{\text{RNS}}$$

$$1 = (1|1|1|1)_{\text{RNS}} \rightarrow -1 = (7|6|4|2)_{\text{RNS}}$$

$$419 = (3|6|4|2)_{\text{RNS}} \rightarrow -419 = (5|1|1|1)_{\text{RNS}}$$



# ***Difficult RNS Operations***

- ▶ Speed of addition and multiplication in RNS is counterbalanced by difficulty of other important arithmetical operations:
  - division
  - sign test
  - magnitude comparison
  - overflow detection
- ▶ Applications of RNS are limited to fields with predominant use of addition and multiplication within the known range (FFT, DSP)

# **Efficiency of RNS Representation**

- ▶ Residues are kept internally as separate binary numbers (bit-fields)

e.g.  $5 = (5|5|0|2)_{\text{RNS}(8|7|5|3)} \rightarrow 10110100010$  (11-bits)

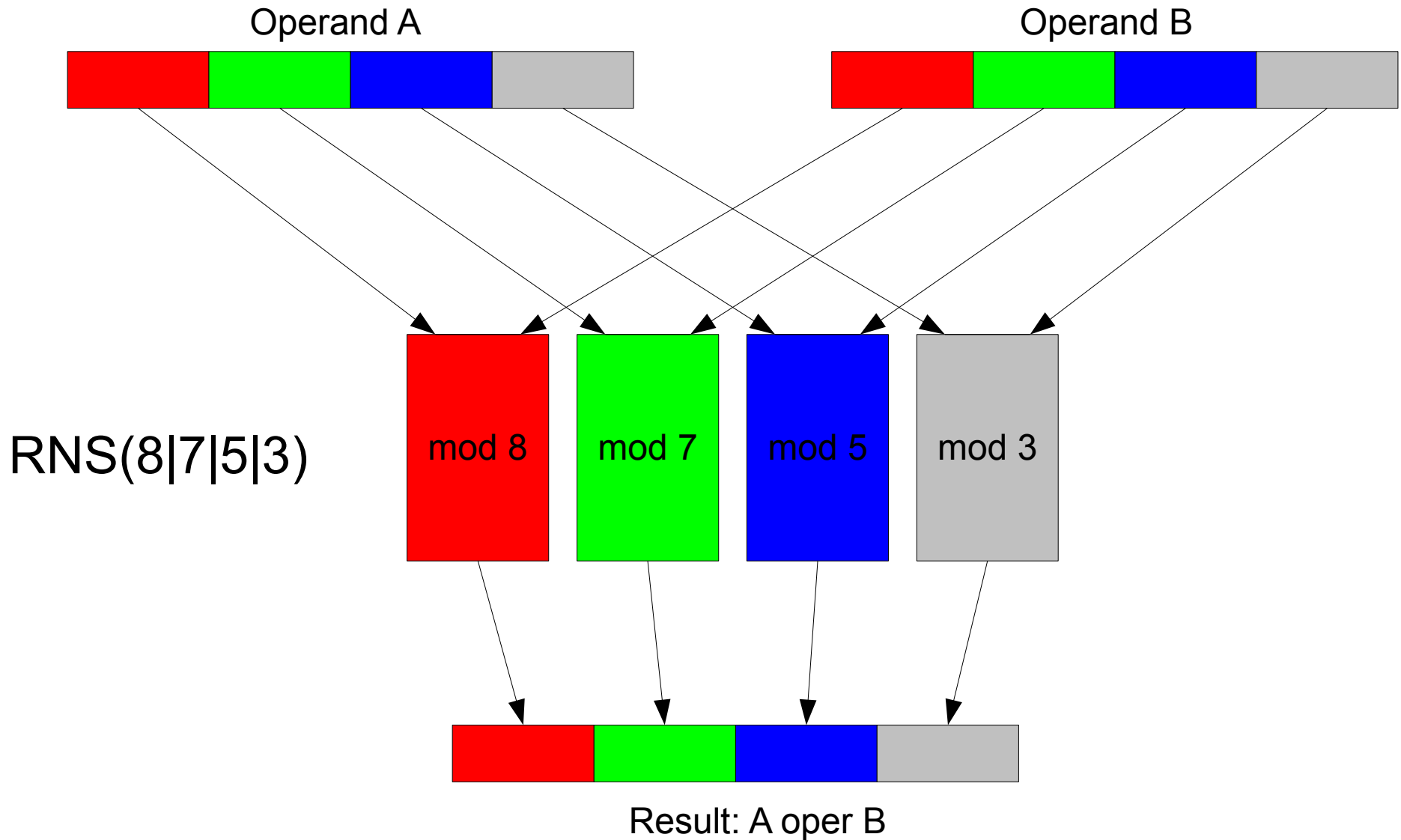
- ▶ How many bits are needed for M-different representations (in NBC)?

$$\log_2 M \rightarrow \log_2 840 = 9.714 \text{ bits}$$

- ▶ Repr. efficiency is the ratio possible RNS representations related to NBC with the same bit-field length  $n$

- $\text{Eff} = M(n) / 2^n \rightarrow 840/2048 = 41\%$

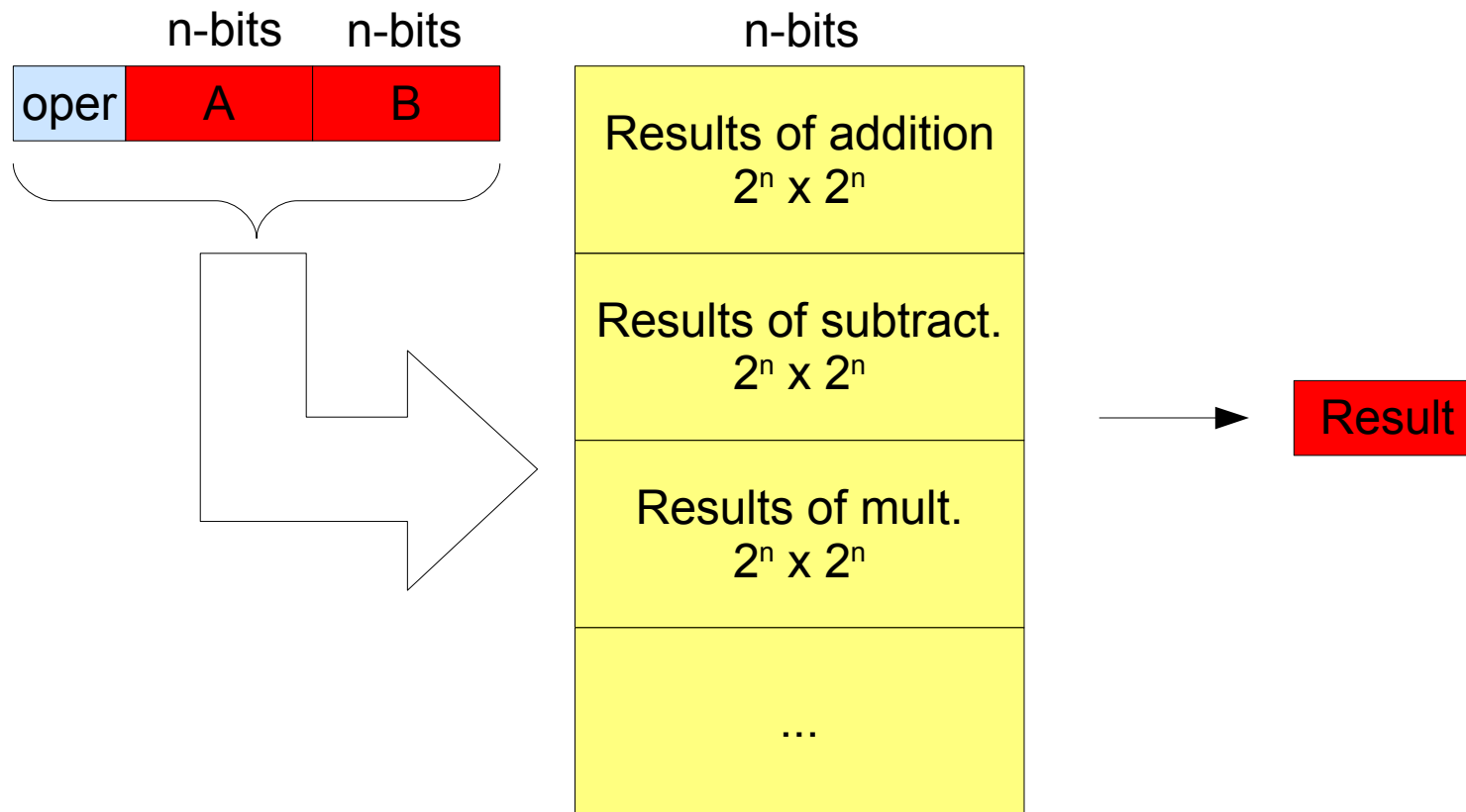
# RNS Arithmetical Unit





# Fast Arithmetics with LUT

- ▶ Small size of operands permits the lookup-table implementation of residue arithm.units



e.g. 4-operations on 4-bits operands  $\rightarrow 4 * (2^4 * 2^4) = 4*256 = 1024$  4bit words (512B)



# Choosing the RNS Moduli

- ▶ Moduli set  $(m_{k-1} \dots m_0)$  affects both
  - representation efficiency
  - complexity of arithmetical units
- ▶ For the chosen range  $M$ :
  - find the moduli: primes (or semi-primes) with product  $\geq M$
  - roughly equal bit-size of moduli should be favored (not an easy task)
  - moduli  $2^n$  and  $2^n-1$  simplify arithmetic units

e.g.  $M=65535 \rightarrow \text{RNS}(13|11|9|7|5|2)$ ,  $\prod m_i=90090$ , 20 bits :-)

# **Choosing the Moduli - Example**

$M = 100\ 000$  (17 bits in NBC)

- ▶ Select consecutive primes

$\text{RNS}(17|13|11|7|5|3|2) \rightarrow \Pi=510510, 23 \text{ bits}$

- ▶ Remove primes to scale the result

$\text{RNS}(17|13|11|7|3|2) \rightarrow \Pi=102102, 20 \text{ bits}$

- ▶ Combine primes – equalize moduli length

$\text{RNS}(26|21|17|11) \rightarrow \Pi=102102, 19 \text{ bits}$

- ▶ Use powers of small primes

$\text{RNS}(15|13|11|2^3|7) \rightarrow \Pi=102102, 18 \text{ bits, max 4-bit field}$

- ▶ Favor  $2^n$  and  $2^n-1$  moduli

$\text{RNS}(2^5|2^5-1|2^4-1|2^3-1) \rightarrow \Pi=104160, 17 \text{ bits, eff}\approx 100\%$

# Conversion to RNS

Calculate the residues  $m_{k-1} \dots m_0$ :

- ▶ From decimal – requires division
- ▶ From binary – requires addition of precomputed values from lookup table and simple division

$$\underbrace{b_{k-1} \dots b_0}_{\text{NBC number}} \bmod m_i = \underbrace{(2^{k-1} b_{k-1} \bmod m_i + \dots + 2^0 b_0 \bmod m_0)}_{\substack{\text{Precomputed} \\ \text{for given RNS}}} \bmod m_i$$

$\underbrace{\hspace{15em}}_{\text{Precomputed for given RNS}}$

$\underbrace{\hspace{15em}}_{\text{requires division of relatively small number - can be implemented with LUT}}$

$\underbrace{\hspace{15em}}_{\text{i-th residue of RNS representation}}$

# Binary to RNS - Example

▶ 164:  $10100100_{NBC} \rightarrow RNS(8|7|5|3) ?$

$$10100100_{NBC} \bmod 8 = 100_{NBC} = 4 \text{ (trivial)}$$

$$10100100_{NBC} \bmod 7 = (2+4+4) \bmod 7 = 3$$

$$10100100_{NBC} \bmod 5 = (3+2+4) \bmod 5 = 4$$

$$10100100_{NBC} \bmod 3 = (2+2+1) \bmod 3 = 2$$

Final division uses small numbers and can also be implemented with LUT

▶ 164:  $10100100_{NBC} \rightarrow (4|3|4|2)$

Lookup table with precomputed partial residues for RNS with moduli (8|7|5|3)

k	$2^k$	$2^k \bmod \dots$			
		8	7	5	3
0	1	1	1	1	1
1	2	2	2	2	2
2	4	4	4	4	1
3	8	0	1	3	2
4	16	0	2	1	1
5	32	0	4	2	2
6	64	0	1	4	1
7	128	0	2	3	2



# Conversion from RNS

- ▶ Any RNS has associated mixed-radix system (MRS)

$$\text{RNS}(m_{k-1} | \dots | m_0) = \text{MRS}(w_{k-1} | \dots | w_0)$$

- ▶ Conversion from RNS to any positional system is possible if the weights  $w_{k-1} \dots w_0$  are known

- ▶ Any RNS number  $(x_{k-1} | \dots | x_0)$  can be expressed as

$$x_{k-1}(1|0|\dots|0) + x_{k-2}(0|1|\dots|0) + \dots + x_1(0|\dots|1|0) + x_0(0|\dots|0|1)$$

- ▶ The weights  $w_{k-1} \dots w_0$  are equal to RNS numbers

$$(1|0|\dots|0), (0|1|\dots|0) \dots (0|\dots|1|0), (0|\dots|0|1)$$



# Conversion from RNS

- ▶ Values of weights for the given RNS can be precomputed and used as constants

$$w_{k-1} = (1|0|\dots|0)_{\text{RNS}}, \dots w_0 = (0|\dots|0|1)_{\text{RNS}}$$

- ▶ All conversions are done modulo M.

e.g. for RNS(8|7|5|3)

$$(1|0|0|0)_{\text{RNS}} = 105$$

$$(0|1|0|0)_{\text{RNS}} = 120$$

$$(0|0|1|0)_{\text{RNS}} = 336$$

$$(0|0|0|1)_{\text{RNS}} = 280$$

$$(3|2|4|2)_{\text{RNS}} = (3*105+2*120+4*336+2*280) \bmod 840 = (2459) \bmod 840 = 779$$

# Conversion from RNS

▶ How to calculate  $(1|0|\dots|0)_{\text{RNS}} \dots (0|\dots|0|1)_{\text{RNS}}$  ?

▶ e.g.  $(1|0|\dots|0)_{\text{RNS}}$  means the number that is dividable by moduli  $m_{k-2} \dots m_0$

▶  $(1|0|\dots|0)_{\text{RNS}} = m_{k-2} * \dots * m_0 * n \quad (n=1,2,\dots,m_{k-1}-1)$

$$\text{e.g. } (1|0|0|0)_{\text{RNS}(8|7|5|3)} = 7*5*3*n$$

▶  $(m_{k-2} * \dots * m_0 * n) \bmod m_{k-1} = 1 \rightarrow n=?$

$$(7*5*3*n) \bmod 8 = 1 \rightarrow n=1$$

$$(1|0|\dots|0)_{\text{RNS}} = 7*5*3 = 105$$

Selection of  $n$  is easy due to very limited range  $[1, \dots, m-1]$